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ADMISSIBLE SLIDES FOR GENERALIZED  
BAUMSLAG–SOLITAR GROUPS

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ABSTRACT. A generalized Baumslag-Solitar group (*GBS* group) is a finitely generated group  $G$  which acts on a tree with all edge and vertex stabilizers infinite cyclic. Any *GBS* group is isomorphic to fundamental group  $\pi_1(\mathbb{A})$  of some labeled graph  $\mathbb{A}$ . Slide is a transformation of labeled graphs. Slides play an important role in isomorphism problem for *GBS* groups. Given an edge  $e$  with label  $\lambda$  and  $\alpha \in \mathbb{Q}$ . In this paper we describe an algorithm that checks if there exists a cycle  $p$  such that after slide  $e$  over  $p$  label  $\lambda$  multiplies by  $\alpha$  or not. If such cycle exists then the algorithm finds one of them.

**Keywords:** isomorphism problem, generalized Baumslag–Solitar group, labeled graph.

## 1. INTRODUCTION

Call a finitely generated group  $G$  a *generalized Baumslag-Solitar group* or a *GBS* group if  $G$  can act on a tree so that the stabilizers of vertices and edges are infinite cyclic groups. By the Bass-Serre theorem,  $G$  is representable as  $\pi_1(\mathbb{A})$ , the fundamental group of a graph of groups  $\mathbb{A}$  (see [1]) with infinite cyclic edge and vertex groups.

Given a *GBS* group  $G$ , we can present the corresponding graph of groups  $\mathbb{A}$  by a labeled graph  $(A, \lambda)$ , where  $A$  is a finite connected graph (with endpoint functions  $\partial_0, \partial_1: E(A) \rightarrow V(A)$ ) and  $\lambda: E(A) \rightarrow \mathbb{Z} \setminus \{0\}$  labels the edges of  $A$ . The label  $\lambda_e$  of an edge  $e$  with the source vertex  $v$  defines an embedding  $\alpha_e: e \rightarrow v^{\lambda_e}$  of the cyclic edge group  $\langle e \rangle$  into the cyclic vertex group  $\langle v \rangle$  (for more details see [5])

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Recently *GBS* groups have been quite actively studied [8], [4], [3]. In particular, the isomorphism problem for *GBS* groups has been discussed: to determine algorithmically when two given labeled graphs define isomorphic *GBS* groups. Despite that, the isomorphism problem is solved only in several special cases [2], [5], [6], but the general solution is not established.

If two labeled graphs  $\mathbb{A}$  and  $\mathbb{B}$  define isomorphic *GBS* groups  $\pi_1(\mathbb{A}) \cong \pi_1(\mathbb{B})$  and  $\pi_1(\mathbb{A})$  is not isomorphic to  $\mathbb{Z}, \mathbb{Z}^2$  or Klein bottle group then there exists a finite sequence of *expansion* and *collapse* (see fig.1) moves connecting  $\mathbb{A}$  and  $\mathbb{B}$  [7]. A labeled graph is called *reduced* if it admits no collapse move (equivalently, the labeled graph contains no edges with distinct endpoints and labels  $\pm 1$ ).

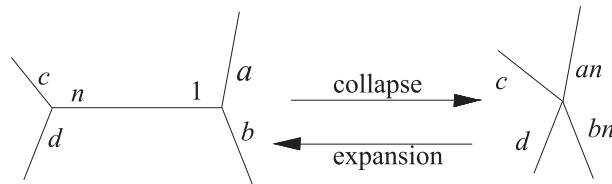


FIG. 1. Expansion and collapse moves.

The following three types of transformations of labeled graphs play an important role for *GBS* groups:

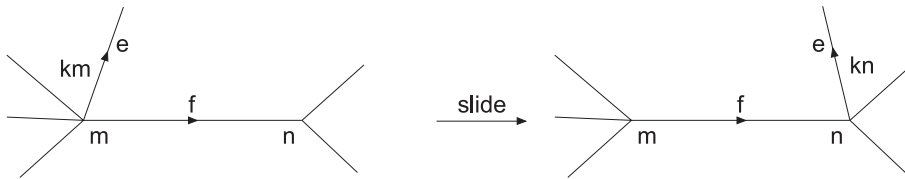


FIG. 2. Slide e/f.

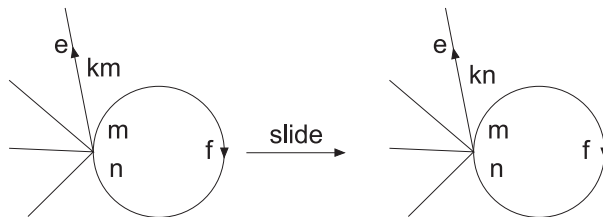


FIG. 3. Slide e/f.

Given a labeled graph  $\mathbb{A}$  (a *GBS* group  $G$ ), denote the set of reduced labeled graphs with the fundamental group isomorphic to  $\pi_1(\mathbb{A})$  (resp.  $G$ ) by  $R(\mathbb{A})$  (resp.  $R(G)$ ).

**Theorem (M. Clay, M. Forester [8])** *Given *GBS* group  $G$  and  $\mathbb{A}, \mathbb{B} \in R(G)$  then labeled graphs  $\mathbb{A}$  and  $\mathbb{B}$  are related by a finite sequence of slides, inductions, and  $\mathcal{A}^{\pm 1}$ -moves, with all intermediate labeled graphs reduced.*

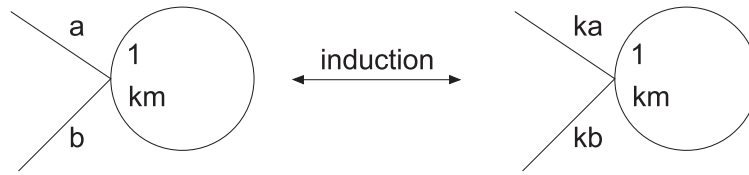


FIG. 4. Induction.

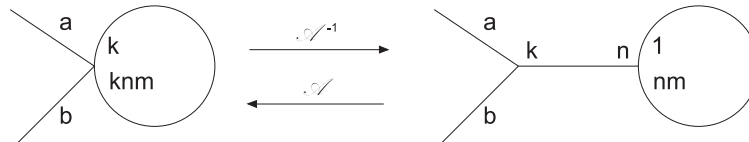


FIG. 5.  $\mathcal{A}$  and  $\mathcal{A}^{-1}$  moves.

An *edge path*  $p = (e_1, e_2, \dots, e_n)$  is a sequence of edges such that  $\partial_0(e_l) = \partial_1(e_{l-1})$  for  $l = 1, \dots, n$ , there  $n = |p|$  is a *length* of  $p$ . A *loop* is an edge  $e$  such that  $\partial_0(e) = \partial_1(e)$ . A *cycle* is an edge path such that  $\partial_0(e_1) = \partial_1(e_n)$ . A geometric edge is a pair of the form  $\{e, \bar{e}\}$ .

Given a labeled graph  $\mathbb{A} = (A, \lambda)$  and an edge path in  $A$   $p = (e_1, e_2, \dots, e_n)$ , put

$$\Delta(p) = \prod_{i=1}^n \frac{\lambda_{\bar{e}_i}}{\lambda_{e_i}}.$$

If the sequences of slides  $e/e_i$  for  $i = 1, \dots, n$  is admissible, then the slide of  $e$  over the path  $p = (e_1, \dots, e_n)$  is called *admissible* and denoted by  $e/p$ . If  $e \in E(A)$  is fixed then path  $p$  is called *admissible* if  $e/p$  is admissible.

Given a graph  $A$ , denote by  $A_0$  the undirected graph obtained from  $A$  by identifying the edges  $e$  and  $\bar{e}$ . If  $T_{A_0}$  some maximal subtree of the  $A_0$ , then denote by  $T_A$  subgraph of  $A$  such that  $(T_A)_0$  coincides with  $T_{A_0}$ .

Suppose that  $\{f_1, \dots, f_k\} = E(A \setminus T_A)$ . Given  $a \in V(A)$ , denote by  $p_i, i = 1, \dots, k$  cycles in  $A$  with endpoints  $a$  such that  $p_i$  one time passes through  $f_i$  and doesn't pass through  $f_j$  or  $\bar{f}_j$  for  $j \neq i$ . Cycles  $p_1, \dots, p_k$  is called a *basic cycles with endpoints a*. A subgroup  $\langle \Delta p_1, \Delta p_2, \dots, \Delta p_k \rangle$  of the multiplicative group  $\mathbb{Q}^*$  is denoted  $\Delta\pi_1(\mathbb{A})$ . The map  $\Delta: \pi_1(\mathbb{A}) \rightarrow \mathbb{Q}^*$  is defined on *GBS* groups and called *modular homomorphism* (see for example [6]).

Suppose that there exist an edge  $e$  and a cycle  $p$  in  $A$  such that after slide  $e/p$  graph  $A$  doesn't change, but the label on  $e$  multiplies by  $\alpha \in \mathbb{Z} \setminus \{0, \pm 1\}$ . In this case slide  $e/p$  can be repeated many times. It turns out that the group  $\pi_1(\mathbb{A})$  can be represented by infinite numbers of reduced labeled graphs and the theorem of M. Clay and M. Forester can not be used to solve the isomorphism problem without restrictions on graph  $A$ .

Known solutions of isomorphism problem was obtained under essential restrictions on  $\Delta\pi_1(\mathbb{A})$  or on the structure of graph  $A$ . These restrictions usually allowed to describe admissible slides and, as a result, to solve the isomorphism problem. In general case to describe an admissible slides is quite difficult problem. In this paper without assumption on labeled graph  $\mathbb{A}$  we try to understand what labels can be written on edge  $e$  after all possible slides of  $e$ .

**Theorem 8** *Given labeled graph  $\mathbb{A}$ ,  $e \in E(A)$  and  $\alpha \in \Delta\pi_1(\mathbb{A})$ . There exists an algorithm that checks if there exists a cycle  $p$  such that slide  $e/p$  is admissible and  $\Delta p = \alpha$  or not. If such cycle exists then the algorithm finds one of them.*

2. DEFINITIONS AND AUXILIARY STATEMENTS

A graph  $A$  is a set  $V(A)$  (vertices), a set  $E(A)$  (edges), maps  $\partial_0, \partial_1: E(A) \rightarrow V(A)$  (endpoints) and a fixed point free involution  $e \rightarrow \bar{e}$  of  $E(A)$  (reversal of orientation) such that  $\partial_i \bar{e} = \partial_{1-i} e$ . Denote by  $A_0$  the undirected graph corresponding to graph  $A$  (we can get  $A_0$  from  $A$  by identifying  $e$  and  $\bar{e}$  for all  $e \in E(A)$ ).

Let  $\{p_1, p_2, \dots, p_r\}$  be the set of prime divisors of  $\prod_{e \in E(A)} \lambda_e$ . If  $\pi_1(\mathbb{C}) \cong \pi_1(\mathbb{A})$  then any label of labeled graph  $\mathbb{C}$  is equal to  $p_1^{\alpha_1} \cdot p_2^{\alpha_2} \cdot \dots \cdot p_r^{\alpha_r}$  for some integers  $\alpha_i \geq 0$  for  $i = 1, 2, \dots, r$ . Therefore it is possible to define a map

$$\overline{\quad}: \Delta\pi_1(\mathbb{A}) \rightarrow \mathbb{Z}^r$$

such that  $\overline{p_1^{\alpha_1} \cdot p_2^{\alpha_2} \cdot \dots \cdot p_r^{\alpha_r}} = (\alpha_1, \alpha_2, \dots, \alpha_r)$ . The multiplication of labels become to the addition of images. Unless otherwise stated, we will use the additive notation for labels and identify labels with elements of  $\mathbb{Z}^r$ . Let  $S$  be a subset of  $\{1, 2, \dots, r\}$ . Denote  $\{1, 2, \dots, r\} \setminus S$  by  $\bar{S}$ .

Given  $v, u \in \mathbb{Z}_{\geq 0}^r$  denote  $i$ -th coordinate of  $v$  by  $[v]_i \in \mathbb{Z}$ . Say that  $v \leq u$  ( $v < u$ ) if  $[v]_i \leq [u]_i$  for all  $i = 1, \dots, r$  (with strict inequality for at least one index). Set  $\mathbb{Z}_{\geq 0}^r$  with binary relation  $\leq$  is a partially ordered set. For  $a, b \in \mathbb{N}$  it's obvious that  $\bar{a} \leq \bar{b}$  ( $\bar{a} < \bar{b}$ ) if and only if  $a|b$  ( $a$  is proper divisor of  $b$ ).

Given an edge path  $p = (e_1, e_2, \dots, e_n)$  in graph  $A$  and  $\Delta p = v^+ - v^-$  where  $v^+, v^- \geq 0$ . Denote  $\lambda_{e_i}$  by  $\alpha_i$  and  $\lambda_{\bar{e}_i}$  by  $\beta_i$ . Define  $\gamma_p$  as the minimal element from  $\mathbb{Z}_{\geq 0}^r$  such that the following inequalities hold

$$v^- + \gamma_p + \sum_{j=1}^i (\beta_j - \alpha_j) - \alpha_{i+1} \geq 0, i = 0, \dots, n - 1.$$

Or equivalently

$$\gamma_p = \max \left\{ \alpha_{i+1} + \sum_{j=1}^i (\alpha_j - \beta_j) - v^-, i = 0, \dots, n - 1 \right\}.$$

**Lemma 1. (admissibility criterion)** *Slide  $e/p$  is admissible if and only if  $\lambda_e \geq \gamma_p + (\Delta p)^-$ .*

*Proof.* The proof is straightforward.

**Proposition 2.** *Given a sequence  $\tilde{v} = \{v_i\}_{i=1}^\infty$  in  $\mathbb{Z}_{\geq 0}^r$  such that  $|v_i|_1 \leq \alpha_i$  (where  $|v|_1$  is the  $L^1$ -norm). There exists a natural number  $L = L(\tilde{\alpha} = \{\alpha_i\}_{i=1}^\infty, r)$  such that there exist  $l_1 < l_2 \leq L$  such that either  $v_{l_1} \geq v_{l_2}$  or  $v_{l_2} \geq v_{l_1}$ .*

*Proof.* Induction on  $r$ . If  $r = 1$ , then  $L = 2$ . If  $r = 2$ , then it is easy to see that  $L$  is not greater than  $\alpha_1$ . Suppose that for dimension  $r - 1$  such  $L$  exists. Note that if  $v_i \not\geq v_1$ , then for some  $k \in \{1, 2, \dots, r\}$  we have  $[v_i]_k < [v_1]_k$ . For  $n \geq 2$  denote the minimal  $i$  such that  $[v_n]_i < [v_1]_i$  by  $\omega(n)$  and the value  $[v_n]_{\omega(n)}$  by  $\zeta(n)$ . It's easy to see that  $0 \leq \zeta(n) \leq \alpha_1 - 1$  and  $\omega(n) \in \{1, 2, \dots, r\}$ . Therefore the number of different pairs  $(\omega(n), \zeta(n))$  is not greater than  $r \cdot \alpha_1$ . If it is impossible to find desirable number  $L = L(\tilde{\alpha}, r)$  for dimension  $r$  then for some pair  $(k, l)$  there exists an infinite subsequence  $\{v_{m_i}\}_{i=1}^\infty$  such that  $\omega(v_{m_i}) = k$  and  $\zeta(v_{m_i}) = l$  for all  $i \geq 1$ .

In this case for any natural number  $L_1$  and any  $l_1 < l_2 \leq L_1$  we have  $v_{m_{l_1}} \not\geq v_{m_{l_2}}$  and  $v_{m_{l_2}} \not\geq v_{m_{l_1}}$ . This is a contradiction with induction assumption, the number of such elements is not greater than  $L(\{\alpha_{m_i}\}_{i=1}^\infty, r - 1)$ . Proposition proved.

**Remark 3.** *If any element of sequence  $\{\alpha_i\}_{i=1}^\infty$  can be computed, then the number  $L = L(\tilde{\alpha}, r)$  can be computed.*

*Proof.* Firstly note that if  $\tilde{\beta} \geq \tilde{\alpha}$  elementwise, then  $L(\tilde{\beta}, r) \geq L(\tilde{\alpha}, r)$ . It follows from the fact that if the sequence  $\tilde{v} = \{v_i\}_{i=1}^\infty$  in  $\mathbb{Z}_{\geq 0}^r$  satisfies conditions  $|v_i|_1 \leq \alpha_i, i \geq 1$  then this sequence satisfies conditions  $|v_i|_1 \leq \beta_i, i \geq 1$ .

Numerate pairs  $(\omega(n), \zeta(n))$  in order of appearing in sequence  $\{(\omega(n), \zeta(n))\}_{n \geq 2}$  such that the first pair  $p_1$  is equal to  $(\omega(2), \zeta(2))$  and the last pair  $p_J = (\omega(n_J), \zeta(n_J))$  has number  $J$  and  $J \leq r \cdot \alpha_1$ . The number  $n_k$  corresponds to the first element  $v_{n_k}$  such that  $[v_{n_k}]_{\omega(n_k)} = \zeta(n_k)$  and  $\omega(n_k)$  is minimal number with condition  $[v_{n_k}]_{\omega(n_k)} < [v_1]_{\omega(n_k)}$ .

Without lost of generality we can assume that  $\tilde{\alpha}$  is nondecreasing sequence. If it's not true, replace  $\{\alpha_i\}_{i=1}^\infty$  by  $\{\max_{1 \leq j \leq i} \alpha_j\}_{i=1}^\infty$  and find upper bound for  $L(\tilde{\alpha}, r)$  using new nondecreasing sequence.

Firstly assume that  $r = 3$ , then  $J < 3 \cdot \alpha_1$ . Since  $\tilde{\alpha}$  is nondecreasing sequence the maximal value of  $L(\tilde{\alpha}, 3)$  can be obtained if number  $n_2$  is the largest possible. Therefore for calculation of upper bound of  $L(\tilde{\alpha}, 3)$  we can get  $n_2 = L(\{\alpha_i\}_{i \geq 2}, 3 - 1) = \alpha_2$ . Then the number of elements  $v_i$  corresponding to the pair  $p_1$  is not greater than  $\alpha_2$ , and for pair  $p_2$  is not greater than  $\alpha_{\alpha_2+1}$ . Thus  $n_3 \leq L(\{\alpha_i\}_{i \geq 1+\alpha_2+\alpha_{\alpha_2+1}}, 2) = \alpha_{1+\alpha_2+\alpha_{\alpha_2+1}}$  and so on. For arbitrary  $i$  we have  $n_i = \alpha_{n_{i-1}+n_{i-1}+\dots+n_1+1}$ . Since  $J \leq 3 \cdot \alpha_1$ ,  $n_{3\alpha_1} = \alpha_{(\sum_{j=1}^{3\alpha_1-1} n_j+1)}$  and the whole number  $L(\{\alpha_i\}_{i \geq 1}, 3)$  is not greater than  $\sum_{j=1}^{3\alpha_1} n_j + 1$ .

The same reasons for arbitrary  $r$ . Since  $n_1 = 2, n_2 \leq L(\{\alpha_i\}_{i \geq 2}, r - 1), n_3 \leq L(\{\alpha_i\}_{i \geq n_2+1}, r - 1), \dots, n_{r \cdot \alpha_1} \leq L(\{\alpha_i\}_{i \geq \sum_{j=1}^{r \cdot \alpha_1-1} n_j+1}, r - 1)$ . Hence  $L(\{\alpha_i\}_{i \geq 1}, r - 1) \leq \sum_{j=1}^{r \cdot \alpha_1} n_j + 1$ . Therefore the upper bound is computable. Remark proved.

**Proposition 4.** *Given a labeled graph  $\mathbb{A}$  and  $e \in E(A)$  with  $\lambda_e = \mu$ . If there exists a cycle  $q$  in  $A$  such that*

1. *slide  $e/q$  is admissible,*
2.  *$\Delta q \geq 0$  and  $\Delta q \neq 0$ ,*

*then there exists computable constant  $C$  and cycle  $q'$  satisfying conditions 1, 2 and*

3.  *$|q'| \leq C = C(\mathbb{A}, \mu)$ .*

*Proof.* Let  $p$  be a cycle of minimal length satisfying conditions 1 и 2. We will find restriction on  $|p|$  using induction on the number of vertices that  $p$  passes through. Number these vertices in order of appearance on  $p$   $v_0, v_1, \dots, v_{k-1}, k < |V(A)|$ .

Let  $k = 0$ , then  $p = (e_1, e_2, \dots, e_{|p|})$ . Denote  $\tau_l = \mu \cdot \Delta(e_1, e_2, \dots, e_l)$  for  $l = 0, 1, \dots, |p|$ . Then  $\tau_{i+1} \leq |\tau_i| + M$  where  $M = \max_{e \in E(A)} |\Delta e|$ . If there exist  $i < j < \lfloor \frac{|p|-1}{2} \rfloor$  such that  $\tau_i \leq \tau_j$  or  $\tau_j \leq \tau_i$  then one of the following holds

1.  $\tau_i \leq \tau_j$  and  $\tau_i \neq \tau_j$  then the path  $e_1 \cdots e_j \cdot \bar{e}_i \cdot \bar{e}_{i-1} \cdots \bar{e}_1$  satisfies conditions 1, 2 and shorter than  $p$ , contradiction,
2.  $\tau_i \geq \tau_j$  and  $\tau_i \neq \tau_j$  then the path  $e_1 \cdots e_i \cdot e_{j+1} \cdots e_{|p|}$  satisfies conditions 1, 2 and shorter than  $p$ , contradiction,
3.  $\tau_i = \tau_j$  then the path  $e_1 \cdots e_i \cdot e_{j+1} \cdots e_{|p|}$  satisfies conditions 1, 2 and shorter than  $p$ , contradiction.

Therefore there no such  $i$  and  $j$ , then by prop.2 we have  $|p| \leq 2L(\{|\mu| + iM\}_{i=1}^\infty, r) + 1 =: L_0(\mathbb{A}, |\mu|)$ .

To be clear we describe the case  $k = 1$ . Let  $p$  is equal to  $p_1 \cdot p_2 \cdots p_T$  where endpoints of  $p_i$  coincides with  $v_0$  and another vertices are equal to  $v_1$ . Denote  $\tau_l = \mu \cdot \Delta(p_1 \cdots p_l)$  for  $l = 0, \dots, T$ , then use above argumentation we can prove that  $|\tau_0| = |\mu|, |\tau_1| \leq |\mu| + 2M + L_0(\mathbb{A}, |\mu| + M), |\tau_2| \leq |\tau_1| + 2M + L_0(\mathbb{A}, |\tau_1| + M), \dots, |\tau_l| \leq |\tau_{l-1}| + 2M + L_0(\mathbb{A}, |\tau_{l-1}| + M)$ . Again, using ideas of the case  $k = 0$  we can obtain restriction  $T \leq 2L(\{|\tau_i|\}_{i=1}^\infty, r) + 1 = L_1(\mathbb{A}, |\mu|)$ .

Step of induction  $s - 1 \rightarrow s$ .

Let  $k = s$  and path  $p$  is equal to  $p_1 \cdot p_2 \cdots p_T$ , where endpoints of  $p_i$  coincides with  $v_0$  and another vertices belong to  $\{v_1, v_2, \dots, v_{s-1}\}$ . Denote  $\tau_l = \mu \cdot \Delta(p_1 \cdots p_l)$  for  $l = 0, \dots, T$ , then use above argumentation we can prove that  $|\tau_0| = |\mu|, |\tau_1| \leq |\mu| + 2M + L_{s-1}(\mathbb{A}, |\mu| + M), |\tau_2| \leq |\tau_1| + 2M + L_{s-1}(\mathbb{A}, |\tau_1| + M), \dots, |\tau_l| \leq |\tau_{l-1}| + 2M + L_{s-1}(\mathbb{A}, |\tau_{l-1}| + M)$ . Again, using ideas of the case  $k = 1$  we can obtain restriction  $T \leq 2L(\{|\tau_i|\}_{i=1}^\infty, r) + 1 = L_s(\mathbb{A}, |\mu|)$ . Proposition proved.

**Example.** Let us construct graph such that the path  $q'$  from prop.4 is very long even with respect to  $|E(A)|$ . We start from graph  $\mathbb{A}_1$  and then construct labeled graphs  $\mathbb{A}_k$  for  $k \in \mathbb{N}$ , see fig.6. Labeled graph  $\mathbb{A}_i$  obtained by attaching of two copies of labeled graph  $\mathbb{A}_{i-1}$  (using vertex  $v_{i-1}$ ) to the labeled graph with two vertices and two edges. Fix  $k \in \mathbb{N}$ , define labeled graph  $\mathbb{B}_k$  by attaching a loop with labels  $p_{k+1}$  and  $p_{k+1} \cdot q$  and an edge  $e$  with labels  $p_1^s$  where  $s = s_1 + s_2 + \dots + s_{k-1} + 1$  (the second label is not important) to the vertex  $v_k$  (see fig.7).

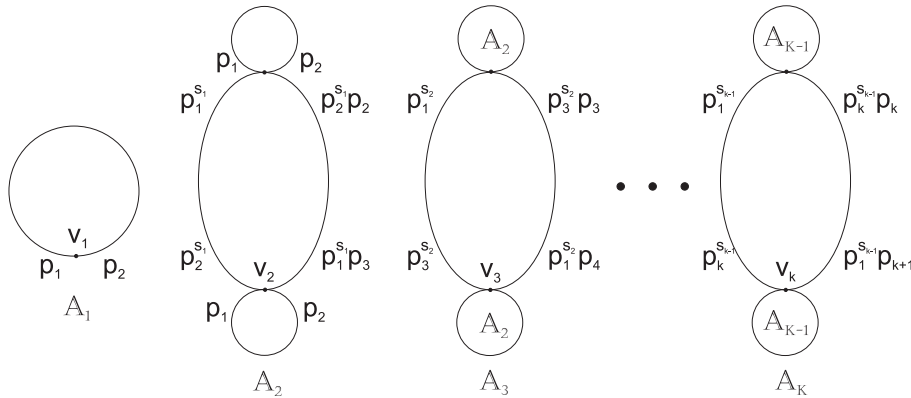


FIG. 6. Labeled graphs  $\mathbb{A}_k$ .

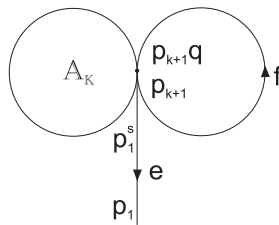


FIG. 7. Labeled graphs  $\mathbb{B}_k$ .

Then graph  $B_k$  has  $2^{k-1} + 1$  vertices and  $2^k + 2^{k-1} - 1$  edges. The main property of  $\mathbb{B}_k$  is following: to change label  $p_1$  to  $p_2$  on  $e$  you need one slide, to change  $p_1 \rightarrow p_3$  requires  $\beta_2 = 2s_1 + 3$  slides, to change  $p_1 \rightarrow p_4$  requires  $\beta_3 = s_2 \cdot \beta_2 + 1 + (s_2 + 1)\beta_2 + 1 = (2s_2 + 1)\beta_2 + 2$  slides, ..., to change  $p_1 \rightarrow p_k + 1$  requires  $\beta_k = (2s_{k-1} + 1)\beta_{k-1} + 2$ . Therefore if path  $p$  satisfies conditions 1 and 2 of prop. 4, then it's length is not less than  $2^k s_1 \cdot s_2 \cdot \dots \cdot s_{k-1}$ . Take  $s_i = 5$  for all  $i$  then  $2 \cdot 10^{k-1}$  is the lower bound on  $|p|$ .

For  $v \in \mathbb{Z}^r$  denote the set  $\{i | [v]_i \neq 0\}$  by  $\text{supp } v$ .

**Proposition 5.** *Given a labeled graph  $\mathbb{A}$  and  $e \in E(A)$  with label  $\lambda_e = \mu$ . Let  $q_1, q_2, \dots, q_k$  be admissible cycles in  $A$  and  $\Delta q_i = v_i \geq 0, i = 1, 2, \dots, k$ . If there exists such admissible cycle  $p$  that*

1.  $\Delta p \geq 0$ ,
  2.  $\text{supp } \Delta p \setminus \bigcup_{i=1}^k \text{supp } v_i \neq \emptyset$ ,
- then there exist computable constant  $C_1$  and admissible cycle  $q$  satisfying 1, 2 and

3.  $|q| \leq C_1 = C_1(\mathbb{A}, \mu, q_1, \dots, q_k)$ .

*Proof.* Numerate coordinates such that  $\bigcup_{i=1}^k \text{supp } v_i = \{1, 2, \dots, l\}$ . Since cycles  $q_1, q_2, \dots, q_k$  are admissible and  $\Delta q_i = v_i \geq 0$ , for any element  $n = (n_1, n_2, \dots, n_l, 0, \dots, 0) \in \mathbb{Z}^r$  we can find constants  $t_1, t_2, \dots, t_k$  such that  $\mu + \sum_{i=1}^k t_i v_i \geq n$ .

Denote by  $\mathbb{A}(q_1, \dots, q_k)$  the labeled graph obtained from  $\mathbb{A}$  as follows: graph  $A(q_1, \dots, q_k)$  coincides with  $A$ , if  $\lambda_f = (\lambda_1, \lambda_2, \dots, \lambda_r)$  the label in  $\mathbb{A}$  then  $\delta_f = (\lambda_{l+1}, \lambda_{l+2}, \dots, \lambda_r)$  is a corresponding label in  $\mathbb{A}(q_1, \dots, q_k)$ .

Since  $p$  satisfies condition 2, we can use prop. 4. Therefore, there is a cycle  $q'$  with  $\Delta q' |_{\{l+1, \dots, r\}} \geq 0$  and  $\neq 0$ . Denote  $\Delta(q')$  by  $v$ . If  $v \geq 0$ , then  $q'$  is a desirable cycle and  $C_1 = C$ , in other case  $v^- \geq 0, \neq 0$ . But  $\text{supp } v^- \subseteq \bigcup_{i=1}^k \text{supp } v_i$ . Since  $|v| \leq C$ ,  $[v^-]_i \leq C$  for  $i = 1, 2, \dots, l$ . Take  $t_1, t_2, \dots, t_k \in \mathbb{Z}$  such that  $\mu + \sum_{i=1}^k t_i v_i \geq v^-$ . Constants  $t_1, t_2, \dots, t_k \in \mathbb{Z}$  bounded depending on  $C$ . Then it's easy to see that  $q = q_1^{t_1} \cdot q_2^{t_2} \cdot \dots \cdot q_k^{t_k} \cdot q'$  is the path we looking for and  $|q|$  bounded depending on  $|q_i|, i = 1, 2, \dots, k$  and  $C$ . Proposition proved.

**Corollary 6.** *Given a labeled graph  $\mathbb{A}$  and  $e \in E(A)$  then either any admissible cycle  $p$  has  $\Delta p \not\geq 0$  or there exists admissible cycle  $q$  and computable constant  $T$  such that*

1.  $\Delta q \geq 0, \neq 0$ ,
2. If  $p$  is admissible and  $\Delta p \geq 0$  then  $\text{supp } \Delta p \subseteq \text{supp } \Delta q$ ,
3.  $|q| \leq T$ .

Such cycle we will call *maximal nonnegative admissible* (or just *maximal*) cycle in labeled graph  $\mathbb{A}$  and denote  $\text{supp } \Delta q$  by  $S(\mathbb{A}, e)$ .

*Proof.* Using proposition 5 several times we can obtain cycles  $q_1, q_2, \dots, q_k$  such that if  $p$  is admissible and  $\Delta p \geq 0$  then  $\text{supp } \Delta p \subseteq \bigcup_{i=1}^k \text{supp } \Delta q_i$ . Denote  $q_1 \cdot q_2 \cdot \dots \cdot q_k$  by  $q$ , then  $\Delta q \geq 0$ , cycle  $q$  is admissible,  $\text{supp } \Delta q = \bigcup_{i=1}^k \text{supp } \Delta q_i$  and  $|q| \leq \sum_{j=2}^l C'_1(\mathbb{A}, \mu, q_1, \dots, q_j) + C = T$ . Corollary proved.

### 3. MAIN RESULTS

Given a labeled graph  $\mathbb{A}$  and  $e \in E(A)$ . Define a graph  $\mathfrak{A} = \mathfrak{A}(\mathbb{A}, e)$  in following way: vertices of graph  $\mathfrak{A}$  is a labeled graphs obtained from  $\mathbb{A}$  by sliding of edge  $e$ , two such labeled graphs  $\mathbb{B}_1$  and  $\mathbb{B}_2$  joined by edge in  $\mathfrak{A}$  if there is an edge  $f$  in  $A$  such that  $\mathbb{B}_1$  is obtained from  $\mathbb{B}_2$  by slide  $e/f$ . It can be established (see prop. 3.10

[5]), that if  $\mathbb{A}$  has no admissible cycle  $p$  such that  $\Delta(p) \geq 0, \neq 0$ , then graph  $\mathfrak{A}$  is finite.

**Lemma 7.** *Given a labeled graph  $\mathbb{A}$  and an edge  $e \in E(A)$ . There exists an algorithm that finds constant  $D$  and admissible nonnegative cycles  $q_1, \dots, q_k$  such that if  $p$  is admissible cycle with  $\Delta p \geq 0$ , then  $\Delta p = \sum_{i=1}^k x_i \Delta q_i$  for some integers  $x_1, \dots, x_k$  and  $|q_i| \leq D$  for  $i = 1, \dots, k$ .*

We will call the set of cycles  $q_1, \dots, q_k$  *basis of admissible nonnegative cycles* (or just *basis*) of  $\mathbb{A}$  with respect to  $e$ .

*Proof.* The corollary 6 implies that if  $p$  is admissible and  $\Delta p \geq 0$ , then  $\text{supp} \Delta p \subseteq S(\mathbb{A}, e)$ . Furthermore  $p$  is admissible in  $\mathbb{A}(q)$  and  $\Delta p|_{\overline{S(\mathbb{A}, e)}} = 0$ . It means that in  $\mathfrak{A} = \mathfrak{A}(\mathbb{A}(q), e)$  cycle  $p$  induce cycle with endpoints  $\mathbb{A}(q)$ . Since labeled graph  $\mathbb{A}(q)$  has no nonnegative admissible cycles, we can see that  $\mathfrak{A}$  is a finite graph and the number of it's vertexes and edges is bounded depending on  $\mathbb{A}$ . Suppose that  $f_1^*, \dots, f_k^*$  are basic cycles in graph  $\mathfrak{A}$  with endpoints  $\mathbb{A}(q)$ , then they induce cycles  $f_1, \dots, f_k$  in  $A$  such that  $|f_i| < 2|E(\mathfrak{A})|$ ,  $0 = \Delta f_i|_{\overline{S(\mathbb{A}, e)}}$  (therefore  $\text{supp} \Delta f_i \subseteq S(\mathbb{A}, e)$ ) and  $\Delta p \in \langle \Delta(f_1), \dots, \Delta(f_k) \rangle_{\mathbb{Z}}$ .

Then for any  $i = 1, \dots, k$  we can find  $y_i \in \mathbb{Z}$  such that cycle  $q_i = q^{y_i} \cdot f_i$  is admissible in  $\mathbb{A}$ . We can calculate  $y_i$  in following way: after admissible slide  $e/q^{y_i}$  the label on edge  $e$  is equal to  $\mu + y_i \cdot \Delta(q)$ . Then by lemma 1 cycle  $f_i$  is admissible if and only if

$$\Delta(f_i^-) + \gamma_{f_i} \leq \mu + y_i \cdot \Delta(q)$$

This inequality is equivalent to the system of two inequalities

$$\Delta(f_i^-)|_{S(\mathbb{A}, e)} + \gamma_{f_i}|_{S(\mathbb{A}, e)} \leq \mu|_{S(\mathbb{A}, e)} + y_i \cdot \Delta(q)|_{S(\mathbb{A}, e)},$$

$$\Delta(f_i^-)|_{\overline{S(\mathbb{A}, e)}} + \gamma_{f_i}|_{\overline{S(\mathbb{A}, e)}} \leq \mu|_{\overline{S(\mathbb{A}, e)}} + y_i \cdot \Delta(q)|_{\overline{S(\mathbb{A}, e)}}.$$

The second inequality holds because  $f_i$  is admissible in  $\mathbb{A}(q)$  for  $i = 1, 2, \dots, k$ . In the first inequality left side is bounded. Therefore the value of  $y_i$  satisfying conditions can be founded.

We have

$$\begin{aligned} \langle \Delta q_1, \dots, \Delta q_k, \Delta q \rangle_{\mathbb{Z}} &= \langle y_1 \Delta q + \Delta f_1, \dots, y_{k-1} \Delta q + \Delta f_k, \Delta q \rangle_{\mathbb{Z}} = \\ &= \langle \Delta f_1, \dots, \Delta f_k, \Delta q \rangle_{\mathbb{Z}} = \langle \Delta f_1, \dots, \Delta f_k \rangle_{\mathbb{Z}}. \end{aligned}$$

The last equality holds because  $q \in \langle \Delta f_1, \dots, \Delta f_k \rangle$  (it's true because  $q$  is admissible in  $\mathbb{A}$  and in  $\mathbb{A}(q)$  too).

The following algorithm finds the basis:

1. Construct  $\mathfrak{A}$ .
2. Find  $f_i^*$  (using a maximal subtree in  $\mathfrak{A}$ ).
3. Induce  $f_1, \dots, f_k$ .
4. Find  $y_i$  using system of inequations.
5. Cycles  $q, q_i = q^{y_i} \cdot f_i, i = 1, \dots, k$  is the basis.

Lemma proved.

**Theorem 8.** *Given a labeled graph  $\mathbb{A}$ , an edge  $e \in E(A)$  and  $v \in \Delta \pi_1(\mathbb{A})$ . Then there exists an algorithm that either finds a path  $p$  such that  $e/p$  is admissible and  $\Delta p = v$  or says that there no such a path.*



*Proof.* Case 1. Let  $v \geq 0$ . If a path  $p$  exists then by lemma 7 we have  $\Delta p = \sum_{i=1}^k x_i \Delta q_i$ . If  $\text{supp } v \not\subseteq S(\mathbb{A}, e)$  then the answer "no" because corollary 6 implies  $\text{supp } \Delta p \subseteq S(\mathbb{A}, e)$ . But if  $\text{supp } v \subseteq S(\mathbb{A}, e)$  then we can find solution of the system

$$v = \sum_{i=1}^k x_i q_i.$$

If there is no solution then the answer "no" (because of lemma 7). Otherwise cycle  $q_1^{x_1} \cdots q_k^{x_k}$  is admissible and desirable.

Case 2. Let  $v \leq 0$  and  $v|_{\overline{S(\mathbb{A}, e)}} = 0$ . Denote by  $\mathbb{B}$  the labeled graph obtained from  $\mathbb{A}$  by replacing of label  $\mu$  by label  $\mu + v$  ( $\mu + v \geq 0$  otherwise the answer is "no"). Initial question is equivalent to the following: to find a cycle  $p_0$  admissible in  $\mathbb{B}$  such that  $\Delta p_0 = -v \geq 0$ . This is the case 1. If the answer "no" then we give the same answer, otherwise  $p = p_0^{-1}$ .

General case. Let  $q$  be a maximal cycle in  $\mathbb{A}$ . Construct a labeled graph  $\mathbb{A}(q)$  and finite graph  $\mathfrak{A} = \mathfrak{A}(\mathbb{A}(q), \mu)$ . Looking for the path in  $\mathfrak{A}$  that connects  $\mathbb{A}(q)$  and labeled graph  $\mathbb{B}$  obtained from  $\mathbb{A}(q)$  by replacing of label  $\mu|_{\overline{S(\mathbb{A}, e)}}$  to the label  $(\mu + v)|_{\overline{S(\mathbb{A}, e)}}$ . If there no such a path (more precisely, there no such vertex  $\mathbb{B}$  in graph  $\mathfrak{A}$ ) then the answer is "no" (since  $p$  is admissible in  $\mathbb{A}$ ,  $p$  have to be admissible in  $\mathbb{A}(q)$  and, therefore,  $p$  induces a path in  $\mathfrak{A}$ ), otherwise we get path  $p^*$ . Let  $p_1$  be a cycle in  $A$  that induces  $p^*$ . Then

$$\Delta p_1|_{\overline{S(\mathbb{A}, e)}} = v|_{\overline{S(\mathbb{A}, e)}}.$$

If path  $p$  exists, then it's induces in  $\mathfrak{A}$  the path that differs from  $p^*$  by some cycle. Therefore either the system

$$\Delta(p_1) - v = - \sum_{i=1}^k t_i \Delta(q_i)$$

is solvable over  $\mathbb{Z}$ , or the answer is "no". Denote the cycle  $q_1^{t_1} \cdots q_k^{t_k} \cdot p_1$  by  $p_2$ . If  $p_2$  is admissible, then  $p = p_2$ .

If  $p_2$  is not admissible, then there is integer  $x$  (see proof of lemma 6) such that cycle  $q^x \cdot p_2$  is admissible. Denote by  $y$  the maximal integer such that cycle  $p_3 = q^x \cdot p_2 \cdot q^{-y}$  is admissible. If  $y = x$ , then  $p = p_3$ .

Let's note, that the path  $p$  exists if and only if there is a path  $q^*$  admissible in labeled graph  $\mathbb{B}$  (obtained from  $\mathbb{A}$  by replacing of label  $\mu$  by label  $\mu + \Delta(p_3)$ ) such that  $\Delta(q^*) = (y - x)\Delta(q) \leq 0$ . This is the case 2. If the answer in case 2 is negative then we give the same answer. Otherwise we have  $p = p_3 \cdot q^*$ . Theorem proved.

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